#### Chapter 10 (Part 2) Operating Systems



# **Chapter Goals**

- Distinguish between fixed and dynamic partitions
- Define and apply partition selection algorithms
- Explain how demand paging creates the virtual memory illusion
- Explain the stages and transitions of the process life cycle
- Explain the processing of various CPU scheduling algorithms

### Partition Memory Management

- Fixed partitions Main memory is divided into a particular number of partitions
- Dynamic partitions Partitions are created to fit the needs of the programs

#### Partition Memory Management



- At any point in time memory is divided into a set of partitions, some empty and some allocated to programs
- Base register A register that holds the beginning address of the current partition
- Bounds register A register that holds the length of the current partition

### Partition Selection Algorithms

Which partition should we allocate to a new program?

- First fit Allocate program to the first partition big enough to hold it
- Best fit Allocated program to the smallest partition big enough to hold it
- Worst fit Allocate program to the largest partition big enough to hold it

- Paged memory technique A memory management technique in which processes are divided into fixed-size pages and stored in memory frames when loaded into memory
  - Frame A fixed-size portion of main memory that holds a process page
  - Page A fixed-size portion of a *process* that is stored into a memory frame
  - Page-map table (PMT) A table used by the operating system to keep track of page/frame relationships

P1 PMT				
Page	Frame			
0	5			
1	12			
2	15			
3	7			
4	22			



#### Figure 10.7 A paged memory management approach



- To produce a physical address, you first look up the page in the PMT to find the frame number in which it is stored
- Then multiply the frame number by the frame size and add the offset to get the physical address

- Demand paging An important extension of paged memory management
  - Not all parts of a program actually have to be in memory at the same time
  - In demand paging, the pages are brought into memory on demand
- Page swap The act of bringing in a page from secondary memory, which often causes another page to be written back to secondary memory

- The demand paging approach gives rise to the idea of virtual memory, the illusion that there are no restrictions on the size of a program
- Too much page swapping, however, is called thrashing and can seriously degrade system performance.

#### **Process Management**

The Process States



#### Figure 10.8 The process life cycle

## **The Process Control Block**

- The operating system must manage a large amount of data for each active process
- Usually that data is stored in a data structure called a process control block (PCB)
- Each state is represented by a list of PCBs, one for each process in that state

## **The Process Control Block**

- Keep in mind that there is only one CPU and therefore only one set of CPU registers
  - These registers contain the values for the currently executing process
- Each time a process is moved to the running state:
  - Register values for the currently running process are stored into its PCB
  - Register values of the new running state are loaded into the CPU
  - This exchange of information is called a context switch

# **CPU Scheduling**

- CPU Scheduling The act of determining which process in the *ready* state should be moved to the *running* state
  - Many processes may be in the ready state
  - Only one process can be in the running state, making progress at any one time
- Which one gets to move from ready to running?

# **CPU Scheduling**

- Nonpreemptive scheduling The currently executing process gives up the CPU voluntarily
- Preemptive scheduling The operating system decides to favor another process, preempting the currently executing process
- Turnaround time The amount of time between when a process arrives in the ready state the first time and when it exits the running state for the last time

# **CPU Scheduling Algorithms**

#### **First-Come, First-Served**

 Processes are moved to the CPU in the order in which they arrive in the running state

#### **Shortest Job Next**

 Process with shortest estimated running time in the ready state is moved into the running state first

#### **Round Robin**

 Each process runs for a specified time slice and moves from the running state to the ready state to await its next turn if not finished

#### **First-Come, First-Served**

Process	Service time
p1	140
p2	75
р3	320
p4	280
p5	125



# **Shortest Job Next**

 Looks at all processes in the ready state and dispatches the one with the smallest service time

C	) 7	5 20	00 34	40 62	20 940
	p2	р5	р1	p4	рЗ

# **Round Robin**

- Distributes the processing time equitably among all ready processes
- The algorithm establishes a particular time slice (or time quantum), which is the amount of time each process receives before being preempted and returned to the ready state to allow another process its turn

## **Round Robin**

Suppose the time slice was 50



## **History of Unix**



System III & V Family

#### Ubuntu



#### www.ubuntu.com



related projects



invite you to participate too!

The Ubuntu community is built on the ideas enshrined in the Ubuntu Philosophy: that

#### **A Little Hands On**





- www.ubuntu.com UBUNTU Official Site
- http://video.google.com/videoplay?docid=-610449081131189823

#### Homework

Read Chapter Ten, Sections 10.3 – 10.4

- Program Assignment #2 Start working on it!!
- Play With Ubuntu

#### **Have A Great Weekend**

